The Volume Operator in Loop Quantum Gravity

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- 0 Motivation
- 1 Construction / Regularization / Implementation on \mathcal{H}_{kin}
- 2 Evaluation of Matrix Elements (Rep'n. Theory / Combinatorics)
- 3 Spectral Properties.
- 4 To Do.

Gauge Invariant 4-Vertex



Gauge Invariance

$$J_1 + J_2 + J_3 + J_4 \stackrel{!}{=} 0$$

Matrix Element

$$\langle j_{12}|\hat{q}_{123}|j_{12}-1\rangle =$$

$$= \frac{1}{\sqrt{(2j_{12}-1)(2j_{12}+1)}} \left[(j_1+j_2+j_{12}+1)(-j_1+j_2+j_{12})(j_1-j_2+j_{12})(j_1+j_2-j_{12}+1) (j_3+j_4+j_{12}+1)(-j_3+j_4+j_{12})(j_3-j_4+j_{12})(j_3+j_4-j_{12}+1) \right]^{\frac{1}{2}}$$

$$= - \langle j_{12}-1|\hat{q}_{123}|j_{12} \rangle$$

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N-Vertex: Simplified Expression for the Matrix Element

 $< \vec{a} | \hat{q}_{IJK} | \vec{a}' > =$

$$= \frac{1}{4} (-1)^{j_{K}+j_{I}+a_{I-1}+a_{K}} (-1)^{a_{I}-a_{I}'} (-1)^{\sum_{n=I+1}^{J-1}j_{n}} (-1)^{-\sum_{p=J+1}^{K-1}j_{p}} \times \\ \times X(j_{I}, j_{J})^{\frac{1}{2}} X(j_{J}, j_{K})^{\frac{1}{2}} \sqrt{(2a_{I}+1)(2a_{I}'+1)} \sqrt{(2a_{J}+1)(2a_{J}'+1)} \times \\ \times \left\{ \begin{array}{c} a_{I-1} & j_{I} & a_{I} \\ 1 & a_{I}' & j_{I}' \end{array} \right\} \left[\prod_{n=I+1}^{J-1} \sqrt{(2a_{n}'+1)(2a_{n}+1)} (-1)^{a_{n-1}'+a_{n-1}+1} \left\{ \begin{array}{c} j_{n} & a_{n-1}' & a_{n}' \\ 1 & a_{n} & a_{n-1} \end{array} \right\} \right] \times \\ \times \left[\prod_{n=J+1}^{K-1} \sqrt{(2a_{n}'+1)(2a_{n}+1)} (-1)^{a_{n-1}'+a_{n-1}+1} \left\{ \begin{array}{c} j_{n} & a_{n-1}' & a_{n}' \\ 1 & a_{n} & a_{n-1} \end{array} \right\} \right] \left\{ \begin{array}{c} a_{K} & j_{K} & a_{K-1} \\ 1 & a_{L}' & j_{K}' \end{array} \right\} \times \\ \times \left[(-1)^{a_{J}'+a_{J-1}'} \left\{ \begin{array}{c} a_{J} & j_{J} & a_{J-1}' \\ 1 & a_{J-1} & j_{J} \end{array} \right\} \left\{ \begin{array}{c} a_{J-1}' & j_{J} & a_{J}' \\ 1 & a_{J} & j_{J} \end{array} \right\} \left\{ \begin{array}{c} a_{J-1} & j_{J} & a_{J}' \\ 1 & a_{J} & j_{J} \end{array} \right\} \\ - (-1)^{a_{J}+a_{J-1}} \left\{ \begin{array}{c} a_{J}' & j_{J} & a_{J-1}' \\ 1 & a_{J-1} & j_{J} \end{array} \right\} \left\{ \begin{array}{c} a_{J-1} & j_{J} & a_{J}' \\ 1 & a_{J} & j_{J} \end{array} \right\} \left\{ \begin{array}{c} a_{J-1} & j_{J} & a_{J}' \\ 1 & a_{J} & j_{J} \end{array} \right\} \right\} \\ \times \\ \times \prod_{n=2}^{I-1} \delta_{a_{n}}a_{n}' \prod_{n=K}^{N} \delta_{a_{n}}a_{n}' \end{array}$$

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4-Vertex

Analytical Insights

• Special Form: only 1 antisymmetric, D-dim tridiagonal matrix, sign factor $\sigma(123)$ only gives overall scaling of the spectrum.

$$\widehat{q}_{123} = \begin{pmatrix} 0 & -q_1 & 0 & \cdots & 0 & 0 & 0 \\ q_1 & 0 & -q_2 & \cdots & 0 & 0 & 0 \\ 0 & q_2 & 0 & \cdots & 0 & 0 & 0 \\ \vdots & \vdots & \vdots & \ddots & \vdots & \vdots & \vdots \\ 0 & 0 & 0 & \cdots & 0 & -q_{D-2} & 0 \\ 0 & 0 & 0 & \cdots & 0 & q_{D-1} & 0 \end{pmatrix}$$

where $q_k = q_k(j_1, j_2, j_3, j_4)$

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Numerical Results.

Histograms for the generic (gauge invariant) 4-vertex

... up to $j_{\rm max} \leq 126/2$. (By 'generic' we mean excluding co-planar edges.)



Oriented Matroids

Motivation from Vectors I

 \mathbb{R}^3 , \mathcal{M} vector config with sorted ground set $E = (\mathbf{e}_1, \dots, \mathbf{e}_5)$.



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- (i) reorientation $\mathbf{e}_k \rightarrow -\mathbf{e}_k$
- (ii) re-labelling



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Oriented Bases $\mathcal{B}(\mathcal{M})$



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- \blacksquare Family $\mathcal{B}(\mathcal{M})$ of sorted bases
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\$\chi_{\mathcal{B}}(S) = {\$\pm \frac{\pm 1}{2} S \in \mathcal{B}\$}\$
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... in our example $\chi_{\mathcal{B}}(B) = \pm \operatorname{sgn}(\det B)$ (if $\chi_{\mathcal{B}}$ chirotope, then also $-\chi_{\mathcal{B}}$, depending of our notion of 'positive' orientation)

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В	123	124	125	134	135	145	234	235	245	345
$\chi_{\mathcal{B},1}(B)$	+	+	+	+	+	+	+	+	+	+
$\chi_{\mathcal{B},2}(B)$	+	+	+	+	+	+	+	+	+	0
$\chi_{\mathcal{B},3}(B)$	+	+	0	+	+	+	+	+	+	0
$\chi_{\mathcal{B},4}(B)$	+	+	+	+	+	+	0	0	0	0

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Our example

is contained in equiv. class 3 (set $\mathbf{e}_K \to -\mathbf{e}_K$ for K=1,3,4,5 and use properties of det)

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Circuits

► $C = \{C \subseteq E : C \text{ min. lin. dep. }\}$ Min. lin. dep. $0 = \sum_{k=1}^{N(c)} \lambda_K \mathbf{e}_K$ $(\mathbf{e}_K \in C, \ \lambda_K \in \mathbb{R})$



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Signed Subsets, $S \subseteq E$

•
$$C = \{C^+, C^-\}$$
 where $C^{\pm} = \{\mathbf{e}_K : \lambda_K \ge 0\}$
• $(-C)^{\pm} = C^{\mp}$. Both, $C, -C$ contained in C



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e₃ e₄ e₄

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- Relative Sign

$$\operatorname{sgn}_{C}(\mathbf{e}_{K}) = \begin{cases} \pm 1 \ \mathbf{e}_{K} \in C^{\pm} \\ 0 \ \mathbf{e}_{K} \notin C \end{cases}$$





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•
$$\underline{C} = \operatorname{supp} C := C^+ \cup C^-$$



Oriented Matroids

Description of vector config \mathcal{M} over ground set E in terms of $\mathcal{B}(\mathcal{M})$ and $\mathcal{C}(\mathcal{M})$ equivalent.

for every $B\in \mathcal{B}$ and for every $e\in E\backslash B$ there is a unique $\pm C\in \mathcal{C}$ such that

 $B \cup \{e\} \subseteq \underline{C}.$

Given two bases $B_1, B_2 \in \mathcal{B}$, $B_1 = (e, b_2, b_3)$, $B_2 = (f, b_2, b_3)$ we have $B_1 \cup \{f\} = B_2 \cup \{e\} \subseteq \underline{C}$ for one $\pm C \in \mathcal{C}$. It holds that

 $\operatorname{sgn}_C(e) \cdot \operatorname{sgn}_C(f) = \chi_{\mathcal{B}}(B_1) \cdot \chi_{\mathcal{B}}(B_2)$

Oriented Matroids

Motivation from Vectors III

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Can convert between the two equivalent descriptions!

Oriented Matroids

For Di-Graphs

The same combinatorics contained in a directed graph:

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Only two realizations of the more general combinatorial concept of an oriented matroid \mathcal{M} of rank 3 over the ground set E in terms of its signed bases $\mathcal{M} = (E, \mathcal{B})$, respectively signed circuits $\mathcal{M} = (E, \mathcal{C})$.

Oriented Matroids Axiomatic Definition: Signed Circuits

A family C of signed subsets of a finite set E is called the set of signed circuits of an oriented matroid $\mathcal{M} = (E, C)$ on E if

(C0) Non-emptiness: $\emptyset \notin C$ (C1) Symmetry: C = -C, that is for every $C \in C$ also its opposite $-C \in C$.

(C2) Incomparability: if $\underline{C_1} \subseteq \underline{C_2}$ then either $C_1 = C_2$ or $C_1 = -C_2$ $\forall \overline{C_1}, \overline{C_2} \in \mathcal{C}.$

(C3) Elimination: For all $C_1, C_2 \in \mathcal{C}$ with $C_1 \neq -C_2$, if $e \in C_1^+ \cap C_2^- \exists C_3 \in \mathcal{C}$ such that $C_3^{\pm} \subseteq (C_1^{\pm} \cup C_2^{\pm}) \setminus \{e\}$.

Equivalent formulation also in terms of $\mathcal{B}(\mathcal{M}).$ Can be extended to infinite ground sets [Bruhn et al.].

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More Difficult: Higher Valence

Sign factor combinatorics for 4-7-valent non-coplanar vertices

N_v	# triples	$\#\vec{\epsilon}(N_v)$ sprinkled	$\#\vec{\epsilon}$ perm.	$\# \vec{\sigma}$ configs	# realizable reor.
			equiv. classes		equiv. classes
3	1	2	1	1	1
4	4	16	3	3	1
5	10	384	4	4	1
6	20	23 808	41	39	4
7	35	3 486 720	706	673	11
8	56	\geq 747 735 880	28 287		135
9	84	?	?		4 381

Numerical Results.

Histograms for each sigma configuration $\vec{\sigma}$ at the (gauge invariant) 5-vertex

... up to $j_{\text{max}} = 25/2$. The blue is for $\vec{\sigma} = (\sigma_{123}, \sigma_{124}, \sigma_{134}, \sigma_{234}) = (2, 0, 0, 0)$, the green for $\vec{\sigma} = (2, 2, 2, 0)$, and the purple for $\vec{\sigma} = (2, 2, 4, 0)$. Each histogram has 512 bins.



 \hat{V} in LQG

Numerical Results

Histograms for the overall generic (gauge invariant) 5-vertex

... up to $j_{\rm max} \leq 25/2$. (By 'generic' we mean excluding co-planar edges.) Each histogram has 512 bins.



Numerical Results

Smallest non-zero eigenvalues λ_{\min} at the (gauge invariant) 5-vertex



Numerical Results

Largest eigenvalues λ_{max} of the (gauge invariant) 5-vertex



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