

## New approaches to resource theories: relative entropic transmitters & epistemic adjointness <sub>Ryszard Paweł Kostecki</sub>

kostecki@fuw.edu.pl

National Quantum Information Center, University of Gdańsk

Institute of Informatics, Faculty of Mathematics, Physics, and Informatics, University of Gdańsk International Center for Theory of Quantum Technologies, University of Gdańsk



I present two new approaches to resource theories: first provides an implementation of Mielnik's idea of nonlinear transmitters, with  $\ell$ -pullbacks of left and right Brègman strongly quasi-nonexpansive operations replacing the CPTP maps, and concrete examples constructed using duality mapping together with the Mazur (resp., Kaczmarz) map on L<sub>p</sub> (resp., Orlicz) spaces over JBW (resp., W<sup>\*</sup>) algebras; second provides a categorical framework for resource theory of epistemic knowledge, with facts (resp., actions) available to a given user encoded in terms of comonad (resp., monad). The first approach provides some special cases of the latter framework.

 $D_{\Psi}, \overleftarrow{\mathfrak{P}}^{D_{\Psi}}, \overrightarrow{\mathfrak{P}}^{D_{\Psi}}, \mathbf{LSQ}(\Psi), \mathbf{RSQ}(\Psi)$  [0]  $\langle f(y) \rangle$  $f(x) + \nabla f(x)^T (y - x)$ 

Presented at: Quantum Information Days, Center for Theoretical Physics, Polish Academy of Sciences, Warszawa, 22.02.2021

 $D_{\ell,\Psi}$  and  $(\ell, \Psi)$ -transmitters [6, 2] • Given  $Z \subseteq int(efd(\Psi))$ , a set *U*, and a bijection  $\ell$ : *U* → *Z*, we define the *Brègman ℓ*-*information* on *U* 

as  $D_{\ell,\Psi}(\phi,\psi) := D_{\Psi}(\ell(\phi),\ell(\psi)) \ \forall \phi,\psi \in U.$ 

- For  $C \subseteq U$ , if  $\ell(C)$  is convex (resp., affine; closed), then *C* will be called  $\ell$ -convex (resp.,  $\ell$ -affine;  $\ell$ -closed).
- For any  $\ell$ -closed  $\ell$ -convex (resp.,  $(\mathfrak{D}^{G}\Psi \circ \ell)$ -closed  $(\mathfrak{D}^{G}\Psi \circ \ell)$ -convex) set C and any  $\psi \in U$ , a *left* (resp., *right*)  $D_{\ell,\Psi}$ -projection is  $\mathfrak{P}_{C}^{D_{\ell,\Psi}}(\psi) := \mathfrak{P}_{\ell(C)}^{D_{\Psi}}(\ell(\psi))$ (resp.,  $\mathfrak{P}_{C}^{D_{\ell,\Psi}}(\psi) := \mathfrak{P}_{\ell(C)}^{D_{\Psi}}(\ell(\psi)).$
- For  $\emptyset \neq W \subseteq U$  and  $T: \ell(W) \to Z, T^{\ell}: \ell^{-1} \circ T \circ \ell: W \to U$

and RSQ-compositional for any strictly increasing, continuous  $\varphi : \mathbb{R}^+ \to \mathbb{R}^+$ , with  $\varphi(0) = 0$ , and  $\lim_{t\to\infty} (t) = \infty$ . This gives us:

2) for any semi-finite JBW-algebra A with a faithful normal semi-finite trace  $\tau$ , the nonassociative  $(L_{1/\gamma}(A,\tau), \|\cdot\|_{1/\gamma})$  spaces [Abdullaev'84, Iochum'84] are uniformly convex and uniformly Fréchet differentiable for any  $\gamma \in ]0,1[$  [Ayupov'86, Iochum'86]. Introducing the nonassociative Mazur map  $\ell_{\gamma} : A_{\star} \ni \phi =$  $|h_{\phi}| \circ s_{\phi} \mapsto |h_{\phi}|^{\gamma} \circ s_{\phi} \in L_{1/\gamma}(A,\tau)$  for  $\phi = \tau(h_{\phi} \circ \cdot),$  $|h_{\phi}| \in A^+, s_{\phi}^2 = \mathbb{I}$ , we obtain the class  $D_{\ell_{\gamma}, \Psi_{\phi}} : A_{\star} \times A_{\star} \rightarrow$  $[0, \infty]$ . Due to isometric isomorphism of  $L_{1/\gamma}(A,\tau)$  for different  $\tau$  [Ayupov–Abdullaev'89],  $D_{\ell_{\gamma}, \Psi_{\phi}}$  do not de-

Let X be a reflexive Banach space (hence,  $X^{**} \cong X$ ), let  $\Psi: X \rightarrow ] - \infty, \infty$ ] be Legendre (:= essentially strictly convex and essentially Gateaux differentiable). Then: 1) Gateaux derivative of  $\Psi$ ,  $\mathfrak{D}^{G}\Psi$  : int(efd( $\Psi$ ))  $\rightarrow$  int(efd( $\Psi^{F}$ )), is a bijection, where: int := topological interior,  $\emptyset \neq$  efd( $\Psi$ ) := { $x \in X | \Psi(x) \neq \infty$ },  $\Psi^{F}(y)$  :=  $\sup_{x \in X} \{ [[x,y]] - \Psi(x) \}, [[x,y]] := y(x) \forall x, y \in X \times X^{*}.$ 2) *Brègman information*, defined as

 $D_{\Psi}(y,x) := \Psi(y) - \Psi(x) - \left[ \left[ y - x, \mathfrak{D}^{\mathsf{G}} \Psi(x) \right] \right]$ for  $x \in int(efd(\Psi))$  and  $\infty$  otherwise, satisfies: (i)  $D_{\Psi}(x,y) \ge 0$ , with = 0 iff x = y; (ii) if  $x \in int(efd(\Psi))$  and  $\emptyset \neq Q \subseteq int(efd(\Psi))$  is convex closed, then  $\operatorname{arginf}_{z \in O} \{ D_{\Psi}(z, x) \}$  is a singleton,  $\{\mathfrak{P}_{O}^{D_{\Psi}}(x)\}$ , called *left*  $D_{\Psi}$ -projection of x onto Q; (iii) left generalised pythagorean theorem:  $D_{\Psi}(y,x) \ge D_{\Psi}(y,\mathfrak{P}_Q^{D_{\Psi}}(x)) + D_{\Psi}(\mathfrak{P}_Q^{D_{\Psi}}(x),x) \ \forall y \in Q,$ with = iff *Q* is affine, characterises  $\mathfrak{P}_{Q}^{D_{\Psi}}$ ; (iv) if  $\emptyset \neq K \subseteq int(efd(\Psi))$  is such that  $\mathfrak{D}^{G}\Psi(K)$  is convex closed, then  $\operatorname{arginf}_{y \in K} \{ D_{\Psi}(x, y) \}$  is a singleton, and its element, the *right*  $D_{\Psi}$ -projection of x onto *K*, satisfies  $\overrightarrow{\mathfrak{P}}_{K}^{D_{\Psi}}(x) = \mathfrak{D}^{G}\Psi^{F} \circ \overleftarrow{\mathfrak{P}}_{\mathfrak{D}^{G}\Psi(K)}^{D_{\Psi}F} \circ \mathfrak{D}^{G}\Psi(x);$ (v) right generalised pythagorean theorem:  $D_{\Psi}(x,y) \ge D_{\Psi}(x, \overrightarrow{\mathfrak{P}}_{K}^{D_{\Psi}}(x)) + D_{\Psi}(\overrightarrow{\mathfrak{P}}_{K}^{D_{\Psi}}(x),y) \ \forall y \in K,$ with = iff  $\mathfrak{D}^{G}\Psi(K)$  is affine, characterises  $\mathfrak{P}_{O}^{D_{\Psi}}$ . 3) Given  $\emptyset \neq M \subseteq int(efd(\Psi))$  and a function  $T: M \rightarrow M$ int(efd( $\Psi$ )), Fix(T) := { $x \in M \mid T(x) = x$ }  $\neq \emptyset$ , while Fix(*T*) is defined as a set of such  $x \in M$  that there exists a sequence  $\{x_n\}_{n \in \mathbb{N}} \subseteq M$  weakly convergent to *x* with  $\lim_{n\to\infty} ||x_n - Tx_n||_X = 0$ . In general,  $\operatorname{Fix}(T) \subseteq \operatorname{Fix}(T)$ . 4) We will call  $T: M \to int(efd(\Psi))$ : (i)  $LSQ(\Psi)$  iff  $D_{\Psi}(x,T(y)) \leq D_{\Psi}(x,y) \ \forall (x,y) \in \operatorname{Fix}(T) \times M \text{ and } (p \in$ Fix(*T*),  $\{y_n\}_{n\in\mathbb{N}}$  is bounded,  $\lim_{n\to\infty} (D_{\Psi}(p,y_n) - p_{\Psi}(p,y_n))$  $D_{\Psi}(p,Ty_n)) = 0) \Rightarrow \lim_{n\to\infty} D_{\Psi}(Ty_n,y_n) = 0;$  (ii)  $RSQ(\Psi)$  iff  $D_{\Psi}(T(x),y) \leq D_{\Psi}(x,y) \quad \forall (x,y) \in$  $M \times \operatorname{Fix}(T)$  and  $(p \in \operatorname{Fix}(T), \lim_{n \to \infty} (D_{\Psi}(y_n, p) -$  $D_{\Psi}(T(y_n),p)) = 0, \quad \{y_n\}_{n \in \mathbb{N}}$  is bounded)  $\Rightarrow$  $\lim_{n\to\infty}(y_n,T(y_n))=0.$ 5) Under some conditions on  $\Psi$  (we will call such  $\Psi$  to be *LSQ-adapted*) one has: (i) if  $\emptyset \neq K \subseteq int(efd(\Psi))$ and  $\{T_1, \ldots, T_n\}$  are LSQ( $\Psi$ ) functions  $K \to K$  such that  $\widehat{F} := \bigcap_{i=1}^{n} \operatorname{Fix}(T_i) \neq \emptyset$  and  $T := T_n \circ \cdots \circ T_1$ , then Fix(T)  $\subseteq F$ , and if Fix(T)  $\neq \emptyset$  then T is LSQ( $\Psi$ ); (ii)

*U* will be called an  $\ell$ -operation (or an  $\ell$ -transmitter). • This gives classes  $LSQ(\ell, \Psi)$  and  $RSQ(\ell, \Psi)$  of  $\ell$ operations. Few other classes are also available. We
will denote  $\widehat{Fix}(T^{\ell}) := \ell^{-1}(\widehat{Fix}(T))$ .

**Resource theories for**  $(\ell, \Psi)$ **-transmitters [1]** 

Given a set *U* (of states), a *resource theory of states* is a triple (P, S, R), where *P* is a submonoid of a monoid of endomorphisms of  $U, \emptyset \neq S \subseteq U$  satisfies  $P(S) \subseteq S$ , and  $R := \{r : U \to \mathbb{R}^+ \mid (r \circ p)(\phi) \leq r(\phi) \forall \phi \in U \forall p \in P\}$ . The elements of *P* (resp., *S*; *R*) are called *free operations* (resp., *free states; resource monotones*). We introduce:

 $1_{\mathrm{R}}^{\mathrm{L}}(\mathcal{T},\operatorname{Fix}(\mathcal{T}),\bigcup_{\phi\in\widehat{\operatorname{Fix}}(\mathcal{T})}\{D_{\ell,\Psi}(\phi,\,\cdot\,)\}) \quad (\operatorname{resp.}, \quad (\mathcal{T},\operatorname{Fix}(\mathcal{T}),$  $\bigcup_{\phi \in \widehat{\operatorname{Fix}}(\mathcal{T})} \{ D_{\ell, \Psi}(\cdot, \phi) \} ) : \text{ if } \emptyset \neq K \subseteq U, \ \mathcal{T} \subseteq \operatorname{LSQ}(\ell, \Psi)$ (resp.,  $\mathcal{T} \subseteq \text{RSQ}(\ell, \Psi)$ ) is a monoid such that  $T^{\ell}: K \to K$  $\forall T^{\ell} \in \mathcal{T}, \ \bigcap_{i=1}^{n} \operatorname{Fix}(T_i) \neq \emptyset \text{ and } \operatorname{Fix}(T_1 \circ \cdots \circ T_n) \neq \emptyset$  $\forall \{T_1^{\ell}, \ldots, T_n^{\ell}\} \subseteq \mathcal{T}$ , then  $D_{\ell, \Psi}(\phi, \cdot)$  (resp.,  $D_{\ell, \Psi}(\cdot, \phi)$ ) is a resource monotone for any  $\phi \in Fix(\mathcal{T})$ ; this holds if Ψ is LSQ-adapted (resp., RSQ-compositional);  $2^{\mathrm{L}}_{\mathrm{R}}(\mathcal{T}, K, \bigcup_{\phi \in K} \{ D_{\ell, \Psi}(\phi, \cdot) \}) \text{ (resp., } (\mathcal{T}, K, \bigcup_{\phi \in K} \{ D_{\ell, \Psi}(\cdot, \phi) \}) \text{):}$ for any  $\ell$ -closed  $\ell$ -convex (resp.,  $(\mathfrak{D}^{G}\Psi \circ \ell)$ -closed  $(\mathfrak{D}^{G} \circ \ell)$ -convex) set  $\varnothing \neq K \subseteq U$ , if  $\mathcal{T}$  is given by the set of all  $\mathfrak{P}_{O}^{D_{\ell,\Psi}}$  (resp.,  $\mathfrak{P}_{O}^{D_{\ell,\Psi}}$ ) with  $\ell$ -closed  $\ell$ -convex (resp.,  $(\mathfrak{D}^{G}\Psi \circ \ell)$ -closed  $(\mathfrak{D}^{G}\Psi \circ \ell)$ -convex) Q such that  $K \subseteq Q$ , with the composition  $\mathfrak{P}_{Q_1}^{D_{\ell,\Psi}} \diamond \mathfrak{P}_{Q_2}^{D_{\ell,\Psi}} := \mathfrak{P}_{Q_1 \cap Q_2}^{D_{\ell,\Psi}}$ (resp.,  $\overrightarrow{\mathfrak{P}}_{Q_1}^{D_{\ell,\Psi}} \diamond \overrightarrow{\mathfrak{P}}_{Q_2}^{D_{\ell,\Psi}} := \overrightarrow{\mathfrak{P}}_{Q_1 \cap Q_2}^{D_{\ell,\Psi}}$ ). Each  $y \in int(efd(\Psi^F))$  defines an *observable* on *U*, given by  $y \circ \ell : U \to \mathbb{R}$ . The (linear) *witnesses* of *S* are defined as the elements of  $\{y \in int(efd(\Psi^F))^+ | [[x,y]] \ge 0 \forall x \in S\}$ . Categories, functors, (co)monads [1] 1. Allowing  $\varnothing$  as object and empty arrows as morphisms, we obtain a category  $lCvx(\ell, \Psi)$  (resp.,  $LSQ_{cvx}(\ell, \Psi)$ ) of  $\ell$ -convex  $\ell$ -closed sets and left  $D_{\ell,\Psi}$ -projections (resp., LSQ( $\ell, \Psi$ ) transmitters), composed by  $\diamond$  (resp.,  $\circ$ ). 2. For any LSQ-adapted  $\Psi$ , there is an embedding functor  $\iota_{\ell,\Psi}^{L}$ :  $\mathbb{ICvx}(\ell,\Psi) \to \mathbb{LSQ}_{Cvx}(\ell,\Psi)$ , and a functor  $\operatorname{Fix}_{\ell,\Psi}^{L}$ :  $LSQ_{CVX}(\ell, \Psi) \rightarrow LCVX(\ell, \Psi)$ , defined by identity on objects and by  $T \mapsto \mathfrak{P}_{Fix(T)}^{D_{\ell,\Psi}}$  on arrows. 3. For  $int(efd(\Psi)) = X = ran(\ell)$ , let Pow(X) be the category of subsets of X and functions between them. A map  $co_{\Psi}^{L}(\cdot)$ , assigning to each  $Y \in Ob(Pow(X))$  the norm closure of a convex hull of Y, and to each  $f \in \operatorname{Arr}(\operatorname{Pow}(X))$  the map  $\mathfrak{P}_O^{D_{\Psi}}$ ,  $Q = \operatorname{co}_{\Psi}^{L}(\operatorname{ran}(f))$ , determines a functor  $\operatorname{co}_{\ell,\Psi}^{\operatorname{L}} : \operatorname{Pow}(\ell^{-1}(X)) \to \operatorname{lCvx}(\ell,\Psi).$ 4. The adjunctions,  $\iota_{\ell,\Psi}^{L} \dashv \operatorname{Fix}_{\ell,\Psi}^{L}$  and  $\operatorname{co}_{\ell,\Psi}^{L} \dashv \operatorname{Frg}_{\operatorname{Set}}$ , where Frg<sub>Set</sub> is a functor forgetting about convexity and topology, equip  $lCvx(\ell, \Psi)$  with a comonad  $co_{\ell \Psi}^{L} \circ$ 

pend on  $\tau$ , only on  $(A, \gamma, \varphi)$ . For  $\varphi(t) = t^{1/\beta - 1}/\beta$ ,  $\beta \in ]0,1[$ , we get  $D_{\gamma,\beta}(\omega, \phi) = (\tau(h_{\omega}))^{\gamma/\beta} + \frac{1-\beta}{\beta}(\tau(h_{\phi}))^{\gamma/\beta} - \frac{1}{\beta}(\tau(h_{\phi}))^{\gamma/\beta - 1}\tau(h_{\omega}^{\gamma} \circ h_{\phi}^{1-\gamma}) \ \forall \omega, \phi \in A_{\star}^{+}.$ 

3) A function  $Y : \mathbb{R} \to \mathbb{R}^+$  is called *Orlicz* iff it is convex,  $Y(0) = 0, Y \not\equiv 0$ , and Y(-u) = Y(u). Any Y and a semifinite W\*-algebra  $\mathcal{N}$  with a faithful normal semi-finite trace  $\tau$  determine a noncommutative Orlicz space  $(L_Y(\mathcal{N}, \tau), \|\cdot\|_Y)$  [Kunze'90]. In [2] we characterise (in terms of conditions on Y and a type of  $\mathcal{N}$ ) strictly convex, uniformly Fréchet differentiable  $L_Y(\mathcal{N}, \tau)$ , satisfying the Radon–Riesz property. By replacing  $[[\cdot, \cdot]]$ with re  $[[\cdot, \cdot]]$ , and introducing the noncommutative Kaczmarz map,  $\ell_Y : \mathcal{N}_* \ni \phi \mapsto u_\phi Y^{-1}(|h_\phi|) \in L_Y(\mathcal{N}, \tau)$ , where  $\phi = \tau(u_\phi |h_\phi| \cdot)$  is a unique polar decomposition, we get  $D_{\ell_Y, \Psi_{\phi}} : \mathcal{N}_* \times \mathcal{N}_* \to [0, \infty]$  (due to isometric isomorphism of  $L_Y(\mathcal{N}, \tau)$  for different  $\tau$  [Ayupov– Chilin–Abdullaev'12], it does not depend on  $\tau$ ).

## Epistemic adjointness [3, 1]

We define an *epistemic inference theory* as a triple (IndInf, E, J), where IndInf is a category of information state spaces as objects and information processings (inductive inferences) as morphisms, J is a monad encoding operations on IndInf, while *E* is a comonad, encoding the range of possible adjunctions  $I \dashv P$  (with a category ExpDes of experimental designs (e.g., spaces of configuration parameters) and experimental procedures,  $I: ExpDes \rightarrow IndInf$  encoding the method of model construction from data, and P: IndInf  $\rightarrow$  ExpDes encoding the criteria of (ideal) experimental verification). • Identifying J with *agent*, E with *coagent*, and (E,J)with a *subject/user*, the *multi-(co)agent* epistemic inference theory is given by (IndInf,  $\{E_i \mid i \in \mathcal{I}\}, \{J_i \mid j \in \mathcal{I}\}$  $\mathcal{J}$ ), and becomes *multi-user* under pairing  $\mathcal{I} = \mathcal{J}$ . • An *intersubjective commensurability* of two subjects/users is given by a groupoid between them, understood internally in the 2-category Comonads × Monads over IndInf, with the (co)lax×(co)lax morphisms as 1-cells, and the choice of lax vs colax (as well as weak vs strong) dependent on the purposes. • Ex.1:  $(1Cvx(\ell,\Psi), co^{L}_{\ell,\Psi} \circ Frg_{Set}, Fix^{L}_{\ell,\Psi} \circ \iota^{L}_{\ell,\Psi})$  is an epistemic information theory with a single user. • Ex.2:  $(\operatorname{Pow}(\mathcal{N}_{\star}), \operatorname{id}_{\operatorname{Pow}(\mathcal{N}_{\star})}, \{\operatorname{Frg}_{\operatorname{Set}} \circ \operatorname{co}_{\ell_{Y}, \Psi_{\omega}}^{L}(\cdot)^{\circ}\}), \text{ with } Y$ and  $\varphi$  varying as in Ex.3) above, is a multi-agent epistemic inference theory. Each agent corresponds to a family of resource theories of states of type  $2^{L}$ ), parametrised by  $\ell_{\rm Y}$ -closed  $\ell_{\rm Y}$ -convex sets of free states. • Ex.3: If IndInf has a terminal object 1, then any agent  $(J, \mu^{J}, \nu^{J})$  determines a monoid  $(M_{I} :=$ Nat $(id_{IndInf}, J), \mu'(\cdot \circ \cdot), \eta')$  of *free operations*, with *resource spaces* (resp., *free resources*) as elements of Ob(IndInf) (resp., { $\sigma_1(1) \in Ob(IndInf) \mid \sigma \in M_I$ }), and functors r: IndInf  $\rightarrow [0, \infty]$  with  $r \circ \sigma_A(A) \leq \sigma_A(A)$  $\forall \sigma \in M_I \; \forall A \in Ob(IndInf) \text{ as resource monotones.}$ 

 $\widehat{\mathfrak{P}}_{Q}^{D_{\Psi}} \in \mathrm{LSQ}(\Psi) \text{ with } \widehat{\mathrm{Fix}}(\widehat{\mathfrak{P}}_{Q}^{D_{\Psi}}) = \mathrm{Fix}(\widehat{\mathfrak{P}}_{Q}^{D_{\Psi}}) = Q.$ 6) (i) Under some conditions on  $\Psi$  (we will call such  $\Psi$  RSQ-compositional): if  $\emptyset \neq K \subseteq X$ ,  $\{T_1, \ldots, T_n\}$  are  $RSQ(\Psi)$  functions  $K \to K$  such that  $\widehat{F} := \bigcap_{i=1}^{n} \widehat{\mathrm{Fix}}(T_i) \neq$   $\emptyset$  and  $T := T_n \circ \cdots \circ T_1$ , then  $\widehat{\mathrm{Fix}}(T) \subseteq \widehat{F}$ , and if  $\widehat{\mathrm{Fix}}(T) \neq \emptyset$  then T is  $RSQ(\Psi)$ ; (ii) under some additional conditions (making  $\Psi$  to be RSQ-adapted) we get  $\widehat{\mathfrak{P}}_{\mathfrak{D}^{\Phi}(M)}^{D_{\Psi}} \in RSQ(\Psi)$  with  $\widehat{\mathrm{Fix}}(\widehat{\mathfrak{P}}_{\mathfrak{D}^{\Phi}(M)}^{D_{\Psi}}) =$  $\operatorname{Fix}(\widehat{\mathfrak{P}}_{\mathfrak{D}^{\Phi}(M)}^{D_{\Psi}}) = \mathfrak{D}^{G}\Psi^{F}(M).$ 

## [0] Key works:

1) Brègman L.M., 1967, *Relaksacionnyĭ metod nakhozhdeniya obšeĭ tochki vypuklykh mnozhestv i ego primenenie dlya resheniya zadach vypuklogo programmirovaniya*, Zh. vychest. mat. mat. fiz. 7, 620–631. mi.mathnet.ru/zvmmf7353 (Engl. transl. 1967, *The relaxation method for finding common points of convex sets and its application to the solution of problems in convex programming*, USSR Comput. Math. Math. Phys. 7, 200–217);

2) Reich S., 1996, A weak convergence theorem for the alternating method with Bregman distances, in: Kartsatos A.G. (ed.), Theory and applications of nonlinear operators of accretive and monotone type, Dekker, New York, pp.313–318;

3) Bauschke H.H., Borwein J.M., Combettes P.L., 2001, *Essential smoothness, essential strict convexity, and Legendre functions in Banach spaces*, Commun. Contemp. Math. **3**, 615–647;

4) Sabach S., 2012, *Iterative methods for solving optimization problems*, Ph.D. thesis, Technion, Haifa.

 $\operatorname{Frg}_{\operatorname{Set}}$  and a monad  $\operatorname{Fix}_{\ell,\Psi}^{\operatorname{L}} \circ \iota_{\ell,\Psi}^{\operatorname{L}}$ .

5. Under some additional conditions,  $\mathfrak{D}^{G}\Psi$  sets functorial equivalence between the above categories and their right versions, inducing the corresponding (comonad, monad) pair on  $rCvx(\ell, \Psi)$ .

## Examples of $D_{\ell,\Psi}$ [2, 1]

1) For a Hilbert space  $\mathcal{H}$ , dim  $\mathcal{H} =: n < \infty$ , the Umegaki relative entropy  $D_1(\rho, \phi) := \operatorname{tr}_{\mathcal{H}}(h_{\rho}(\log h_{\rho} - \log h_{\phi}))$ equals  $D_{\ell=\operatorname{id}, \Psi=\Phi\circ\Lambda}(\rho, \phi)$ , where  $\psi = \operatorname{tr}_{\mathcal{H}}(h_{\psi} \cdot) \in \mathfrak{B}(\mathcal{H})_{\star}$ ,  $\Lambda$  is a nonincreasing rearrangement of eigenvalues, while  $\Phi(x) := \sum_{i=1}^{n} (x_i \log(x_i) - x_i)$  for  $x \ge 0$  and  $\infty$  otherwise [Bauschke–Borwein'97]. So, Lüders' and quantum Jeffrey's rules [4], as well as a partial trace [5], being derived from special cases of  $\mathfrak{P}^{D_1}$ , belong to  $\mathfrak{P}^{D_{\ell},\Psi}$ . If  $(X, \|\cdot\|_X)$  is strictly convex, uniformly Fréchet differentiable, and satisfies the Radon–Riesz property, then  $\Psi_{\varphi} := \int_{0}^{\|\cdot\|_X} \mathrm{d}t \varphi(t) : X \to \mathbb{R}^+$  is Legendre, LSQ-adapted, [1] Kostecki R.P., 2021, Categories of Brègman operations and epistemic (co)monads, arXiv:2103.07810.

[2] Kostecki R.P., 2021, *Postquantum Brègman relative entropies*, arXiv:1710.01837 (version 3 in preparation).

[3] Kostecki R.P., 2016, *Towards (post)quantum information relativity*, Perimeter Institute for Theoretical Physics, Waterloo. PIRSA:16050021.

[4] Hellmann F., Kamiński W., Kostecki R.P., 2016, *Quantum collapse rules from the maximum relative entropy principle*, New J. Phys. 18, 013022. arXiv:1407.7766.
[5] Munk-Nielsen M.I., 2015, *Quantum measurements from entropic projections*, M.Sc. thesis, University of Waterloo and Perimeter Institute of Theoretical Physics, Waterloo. https://www.fuw.edu.pl/~kostecki/morten\_essay.pdf.
[6] Kostecki R.P., 2011, *The general form of γ-family of quantum relative entropies*, Open Sys. Inf. Dyn. 18, 191–221. arXiv:1106.2225.

This research was supported by 2015/18/E/ST2/00327 grant of National Science Centre, MAB/2018/5 grant of Foundation of Polish Science, and by Perimeter Institute for Theoretical Physics. Research at Perimeter Institute is supported by the Government of Canada through Industry Canada and by the Province of Ontario through the Ministry of Research and Innovation.