

How to Cope with the Curse of Dimensionality?

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Curse of Dimensionality

ε error demand

d the number of variables

$n(\varepsilon, d)$ the minimal cost

Many problems suffer from
the *curse of dimensionality*

$$n(\varepsilon, d) \geq c (1 + C)^d$$

for all $d = 1, 2, \dots$ with $c, C > 0$.

Multivariate Integration

For $f \in F_d$ we want to approximate

$$I_d(f) := \int_{[0,1]^d} f(t) dt \approx A_n(f)$$

- **Algorithms:**

$A_n(f) = \phi(f(x_1), f(x_2), \dots, f(x_n))$ with $x_j \in [0, 1]^d$

- **Minimal Worst Case Error:**

$$e(n, d) = \inf_{A_n} \sup_{\|f\|_{F_d} \leq 1} |I_d(f) - A_n(f)|$$

- **Information Worst Case Complexity:**

$$n(\varepsilon, d) = \min\{n \mid e(n, d) \leq \varepsilon\}$$

Multivariate Integration for Smooth Functions

$$L = \{L_d\} \quad L_d > 0$$

$$F_d = C_d^r(L) := \{f : [0, 1]^d \rightarrow \mathbb{R} : \|f\|_{\max} \leq 1, \|D^\alpha f\|_{\max} \leq L_d \forall |\alpha| \in [1, r]\}$$

Bakhvalov [1959]

$$n(\varepsilon, d) = \Theta(\varepsilon^{-d/r})$$

but factors in the Θ -notation depend on d and r . Curse?

Sukharev [1979]: **The curse holds for $r = 1$ and $L_d \equiv 1$.**

Otherwise, curse?

Multivariate Integration for Smooth Functions

$$C_d^r(L) := \{f : [0, 1]^d \rightarrow \mathbb{R} : |f(x)| \leq 1, \quad |D^\alpha f(x)| \leq L_d \ \forall \alpha \in [1, r]\}$$

What are necessary and sufficient conditions for $\{L_d\}$ to have the curse of dimensionality for multivariate integration?

Theorem (Hinrichs, Novak, Ullrich, W [2012])

The curse holds for $C_d^r(L)$ iff $\liminf_{d \rightarrow \infty} L_d \sqrt{d} > 0$

Multivariate Integration for Korobov Spaces

$$r = \{r_j\} \quad \text{with} \quad 1 \leq r_1 \leq r_2 \leq \dots$$

H_{r_j} : 1-periodic $f : [0, 1] \rightarrow \mathbb{C}$, $f^{(r_j-1)}$ abs. cont, $f^{(r_j)} \in L_2$

$$\|f\|_{H_{r_j}}^2 = \left| \int_0^1 f(t) dt \right|^2 + \int_0^1 |f^{(r_j)}(t)|^2 dt$$

For $d \geq 1$,

$$F_d = H_{d,r} = H_{r_1} \otimes H_{r_2} \otimes \dots \otimes H_{r_d}$$

Usually, it is assumed that $r_j \equiv r$

Theorem

Let $r_j \equiv r$. Then there exists $c_r, C_r > 0$ such that

$$n(\varepsilon, d) > c_r (1 + C_r)^d$$

Based on Hickernell+W [2001] and Novak+W[2001], see also Sloan+W[2001]

**Multivariate integration for Korobov space
with arbitrarily smooth functions
suffers from the curse of dimensionality**

How to cope with the curse of dimensionality

- switch to spaces of increased smoothness with respect to successive variables
- switch to weighted spaces, i.e., groups of variables are of varying importance
- switch to a more lenient setting, i.e, from the worst case setting to the randomized or average case setting

Increasing Smoothness

Multivariate integration for Korobov spaces in the worst case setting with $r_1 \leq r_2 \leq \dots$.

But we now allow to increase r_j

Let

$$R := \limsup_{k \rightarrow \infty} \frac{\ln k}{r_k}$$

Theorem

If $R < 2 \ln 2\pi$ then

- **no curse**
- $n(\varepsilon, d) \leq C \varepsilon^{-p(1+p/2)}$ with $p := \max(r_1^{-1}, R/\ln 2\pi) < 2$,
i.e., strong polynomial tractability

Based on Papageorgiou+W [09], Kuo, Wasilkowski+W[09]

Weighted Spaces

Major research activities in last 10 years...

In particular, for $r_j \equiv r$ and $\gamma = \{\gamma_j\}$, redefine H_{r_j, γ_j} with

$$\|f\|_{H_{r_j, \gamma_j}}^2 = \left| \int_0^1 f(t) dt \right|^2 + \frac{1}{\gamma_j} \int_0^1 |f^{(r_j)}(t)|^2 dt$$

For $d \geq 1$,

$$H_{d,r} = H_{r_1, \gamma_1} \otimes H_{r_2, \gamma_2} \otimes \cdots \otimes H_{r_d, \gamma_d}$$

Theorem

- Gnewuch+W[08]

$$\lim_{d \rightarrow \infty} \frac{\sum_{j=1}^d \gamma_j}{d} = 0 \quad \text{iff} \quad \text{no curse,}$$
- Hickernell+W[01]

$$\limsup_{d \rightarrow \infty} \frac{\sum_{j=1}^d \gamma_j}{\ln d} < \infty \quad \text{iff} \quad \text{polynomial tractability,}$$

i.e., $n(\varepsilon, d) \leq C d^q \varepsilon^{-p}$
- Hickernell+W[01]

$$\sum_{j=1}^{\infty} \gamma_j < \infty \quad \text{iff} \quad \text{strong polynomial tractability,}$$

i.e., $n(\varepsilon, d) \leq C \varepsilon^{-p}$

More Lenient Settings

From Worst Case Setting to

- Randomized Setting
- Average Case Setting

Average Case Setting \leq Randomized Setting

Randomized Setting

- **Algorithms:**

$A_{n,\omega}(f) = \phi_\omega(f(x_{1,\omega}), f(x_{2,\omega}), \dots, f(x_{n(\omega),\omega}))$ for a random ω

- **Minimal Randomized Error:**

$$e(n, d) = \inf_{A_n} \sup_{\|f\|_{H_{d,r}} \leq 1} [\mathbb{E} |I_d(f) - A_{n,\omega}(f)|^2]^{1/2}$$

- **Information Randomized Complexity:**

$$n(\varepsilon, d) = \min\{n \mid e(n, d) \leq \varepsilon\}$$

Monte Carlo Algorithm

$$A_{n,\omega}(f) = \frac{1}{n} \sum_{j=1}^n f(x_{j,\omega})$$

with

$x_{j,\omega}$ iid with uniform distribution over $[0, 1]^d$

Sloan+W[01] for Korobov spaces, obvious for $C_d^r(L)$ spaces

- $n(\varepsilon, d) \leq \varepsilon^{-2}$
- no curse and strong polynomial tractability

Conclusions

- Many multivariate problems suffer from the curse of dimensionality in the worst case setting
- We may sometimes break the curse of dimensionality by
 - switching to spaces of increased smoothness with respect to successive variables
 - switching to weighted spaces, i.e., groups of variables are of varying importance
 - switching to a more lenient setting, i.e, from the worst case setting to the randomized or average case setting

Book

More can be found in

Tractability of Multivariate Problems

Erich Novak and Henryk Woźniakowski

- Volume I: Linear Information (2008)
- Volume II: Standard Information for Functionals (2010)
- Volume III: Standard Information for Operators (2012?)